

## • Summary:

**Arithmetic:**  $1 + 2 + 3 + \dots + N$  operations =  $\Theta(N^2)$

$1 + 4 + 7 + \dots + N$  operations =  $\Theta(N^2)$   
 (next number is +c larger)

**Geometric:**  $1 + 2 + 4 + \dots + N$  operations =  $\Theta(N)$

$1 + 8 + 64 + \dots + N$  operations =  $\Theta(N)$   
 (next number is  $\times c$  larger)

## • Analyzing recursive function runtimes:

### 5 STEPS:

① How many **operations** in this function?

② Draw a **tree**! Draw the first function call as a node like so:

③ How many **recursive calls** in this function? Draw as child node as so:



④ Repeat 1-3 on recursive calls until notice pattern

⑤ Label recursive calls by **level** (distance away from first function call) - see example

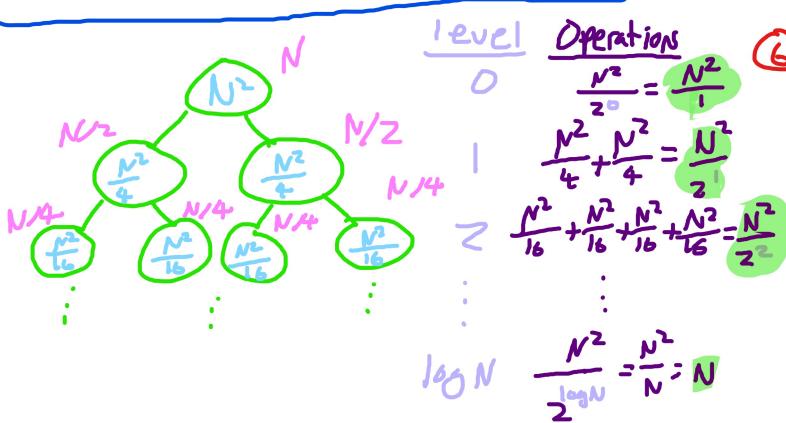
⑥ Add up number of **operations** per level. Then add up operations at each level for each level.

### EXAMPLE:

```
public static void weirdCount(N) {
    if (N <= 1) { return 1; }
    int sum = 0;
    for (int x=0; x < N; x++) {
        for (int y=0; y < N; y++) {
            sum += 1;
        }
    }
    return sum + weirdCount(N/2)
        + weirdCount(N/2);
}
```

- ① In our first function `weirdCount(N)`, there are  $\Theta(N^2)$  operations
- ② Create tree!
- ③ Two recursive calls: `weirdCount(N/2)` and `weirdCount(N/2)`
- ④ Notice that each call to `weirdCount` has two child nodes (recursive calls) with half the input.
- ⑤ We realize that there are  $\log N$  levels because the input  $N$  decreases by half each level;  $N$  halves  $\log N$  times until it reaches 1 (base case=leaf nodes; no more recursive calls).

- ⑥ Notice how we see a pattern in the operations at each level &  $\frac{N^2}{2^0} = \frac{N^2}{1}$ . Now we add up all the operations in the "Operations" column.



$$\begin{aligned}
 & \frac{N^2}{1} + \frac{N^2}{2} + \frac{N^2}{4} + \dots + N \\
 &= N \underbrace{(N + \frac{N}{2} + \frac{N}{4} + \dots + 1)}_{\text{Geometric Sum!}} = N \cdot \Theta(N) = \Theta(N^2)
 \end{aligned}$$